

CSE 2550: Blockchain Technology I

Lecture 14 **Micropayments**

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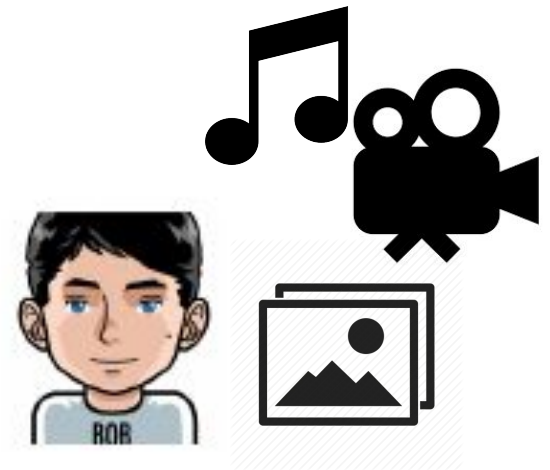
UConn - Fall 2023

Outline

- Motivation.
- On the use of payment channels/networks for micropayments.
- Probabilistic micropayments.
 - Centralized schemes.
 - Decentralized schemes.
 - MICROPAY.
 - MicroCash.



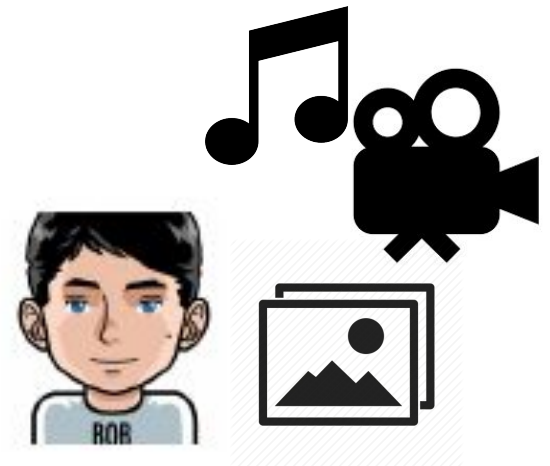
Customer



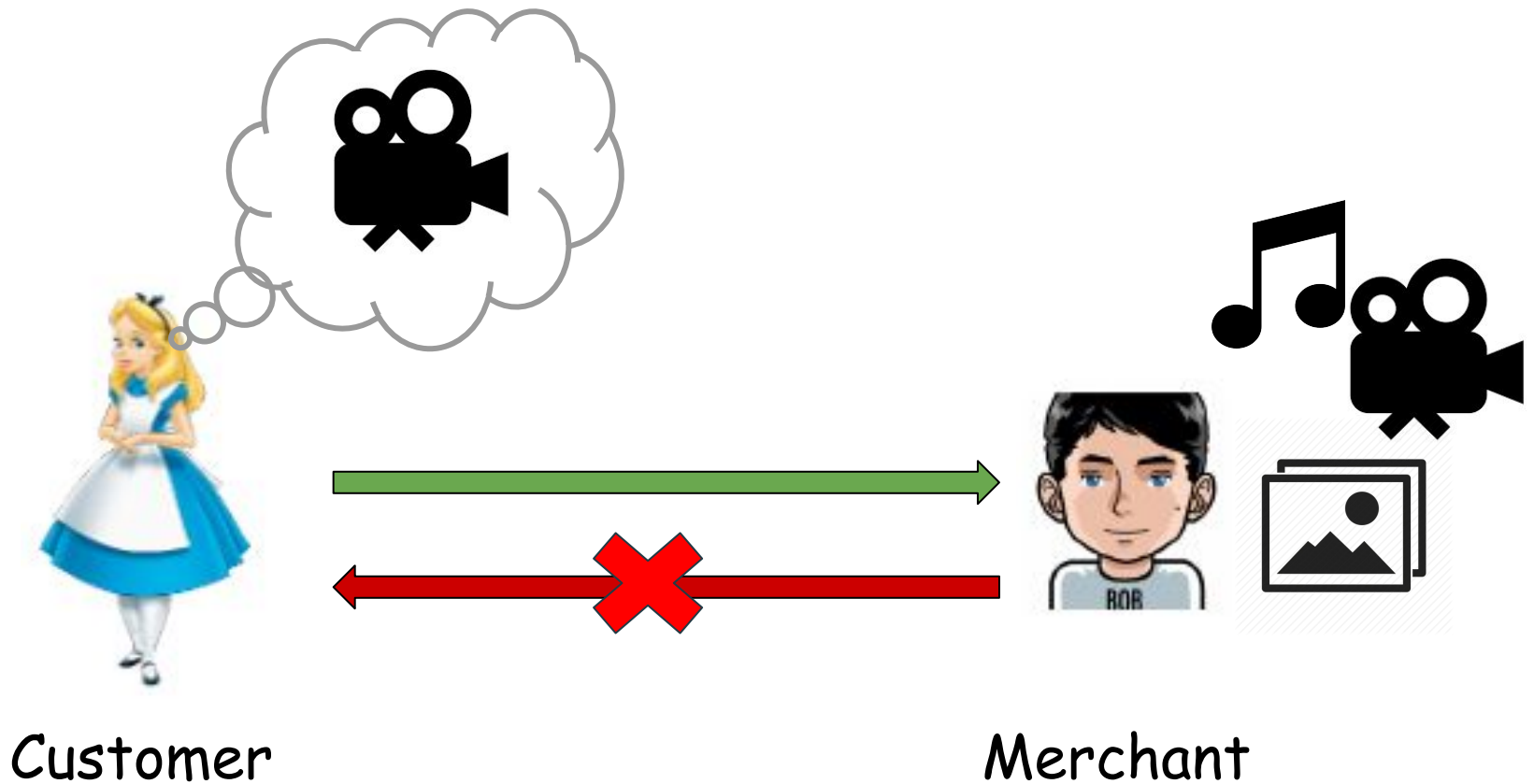
Merchant



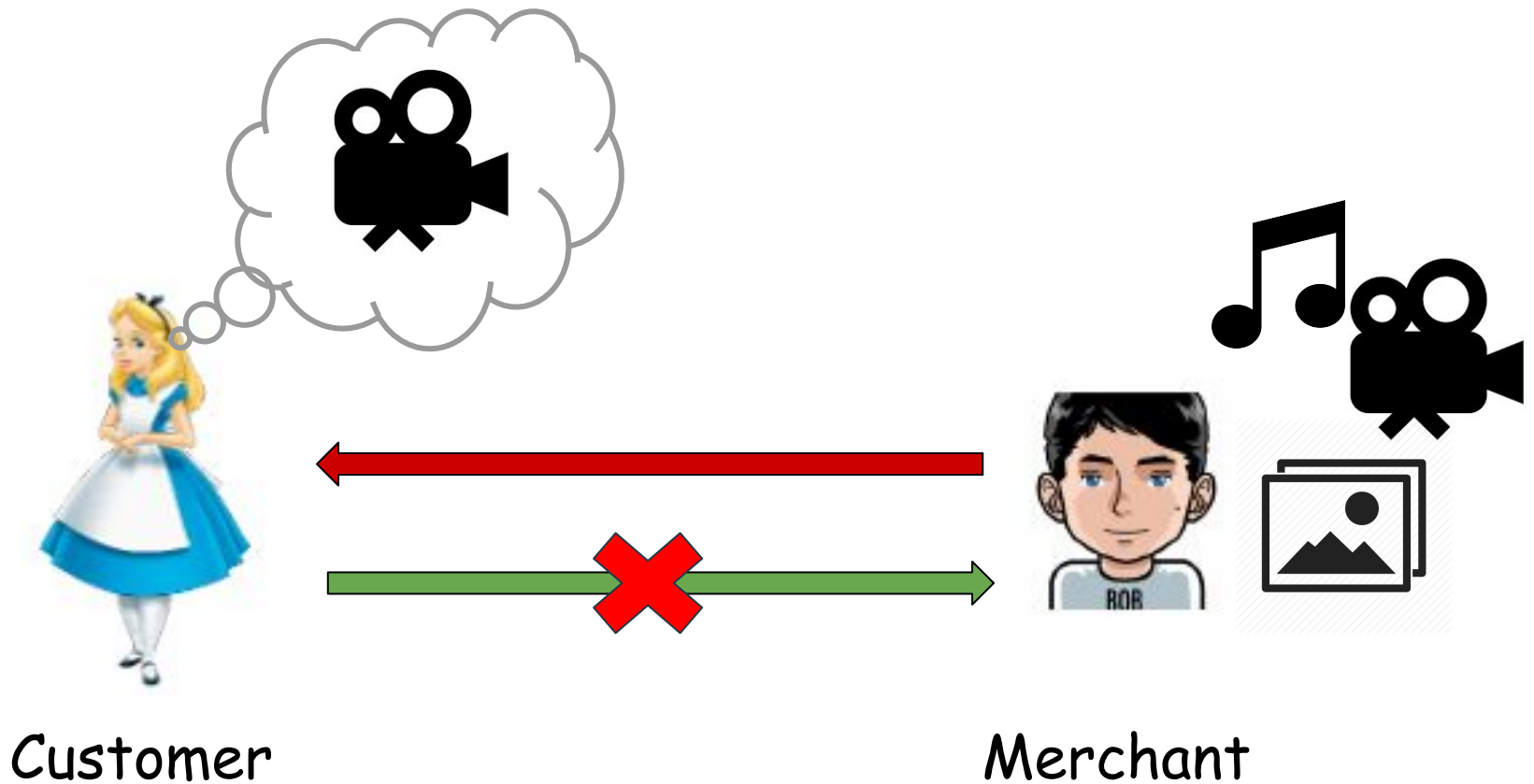
Customer



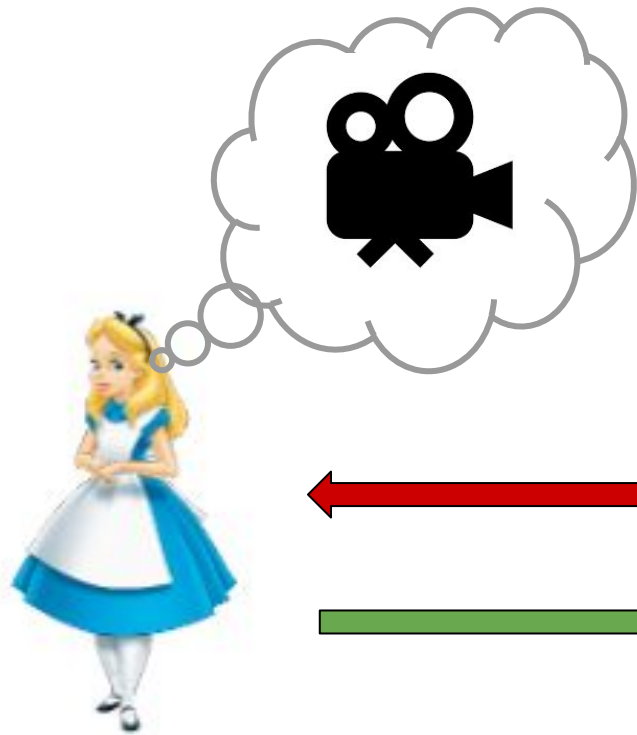
Merchant



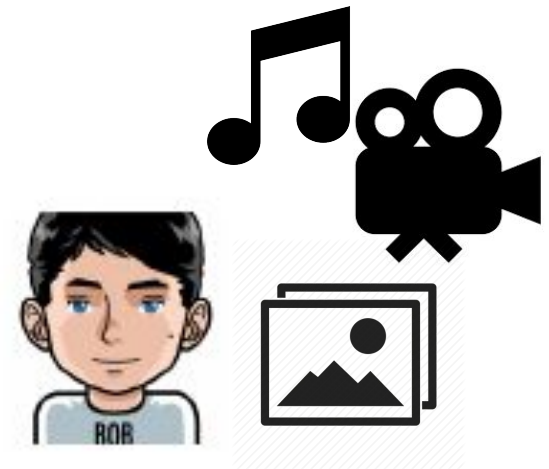
The Merchant could fail to provide the service but keep the customer's money



The Customer could fail to pay after obtaining the service



Customer



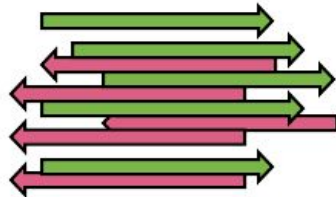
Merchant



I did not like this movie, I just watched the first 30 min!

Micropayments

I am not sure I will like the movie, I will pay per minute I watch!



OK!



Total = \$1
Fees = \$10

We worked a lot!!!!

➡ Micropayment ← Video frame

- Payments of a micro value, i.e., pennies or fractions of pennies.
- Several applications, e.g., ad-free web, online gaming, etc.
 - Used extensively in blockchain-based distributed services.
 - Main motivation is the impossibility of fair service-payment exchange.
 - And flexibility!

Challenges

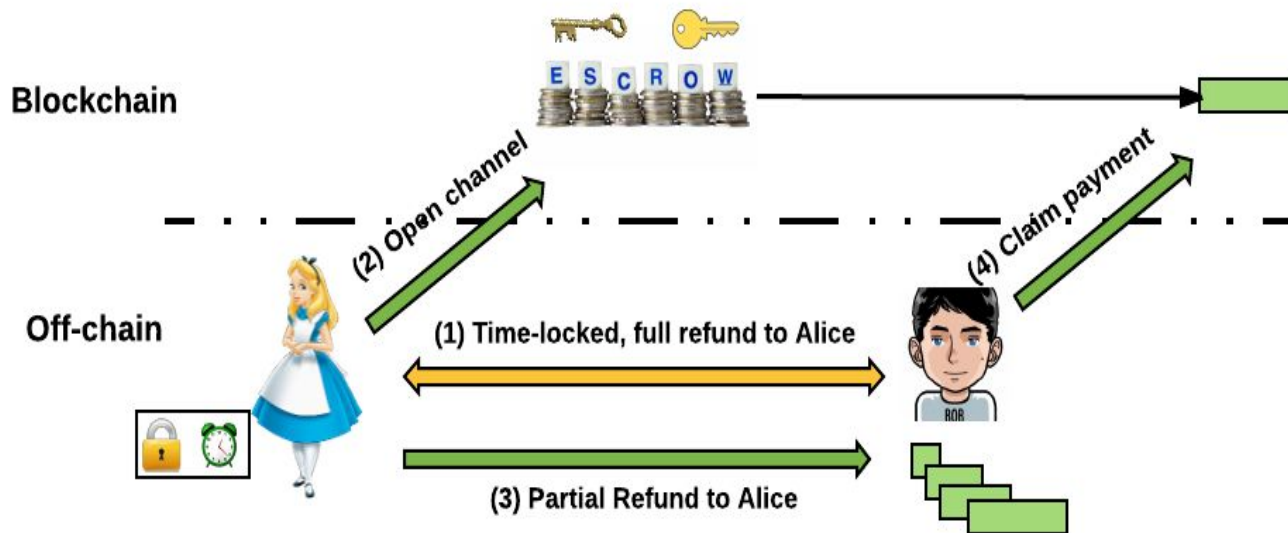
- Produce a huge number of small-value transactions.
 - Overwhelm the system.
 - Explode the payment log.
 - Blockchain storage footprint.
 - Do not scale for large demands or large number of users.
 - What is the throughput of Bitcoin in tx/sec?!
 - High transaction fees.
 - Each transaction must pay a fee.
 - This fee may exceed the payment value itself.

**Aggregate the small payments into
a few larger ones!**



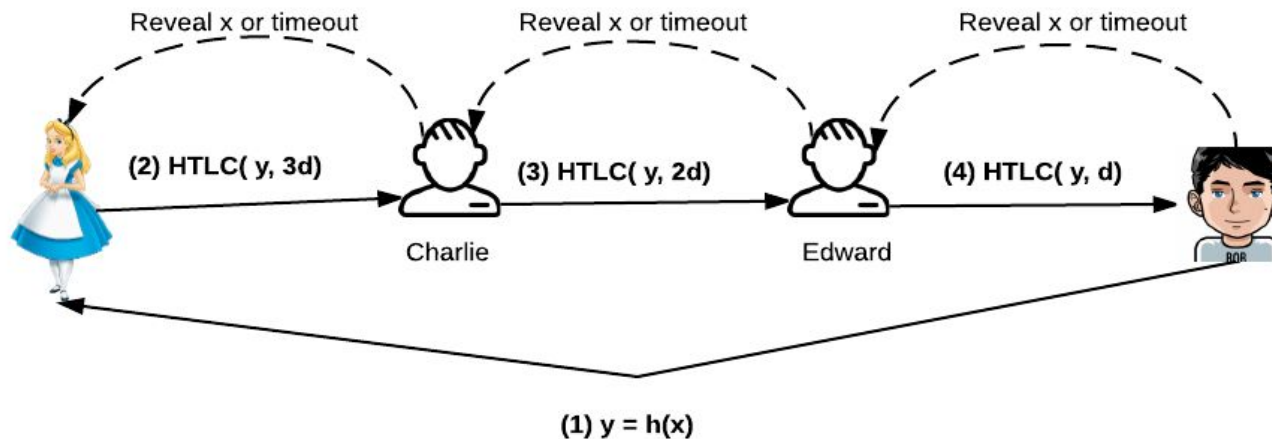
Micropayment Channels (or payment channels)

- Process most transactions off-chain, only channel opening and closing transactions will be on-chain.
- A channel allows exchanging payments between only two parties.



Micropayment Networks (or payment networks)

- Payment networks allow paying several parties.
 - E.g., Lightning networks.
 - Alice can pay Bob as long as there is a payment path between them: a multi-hop or set of consecutive payment channels.
 - Core component: HTLC (Hash Time-Lock Contract).

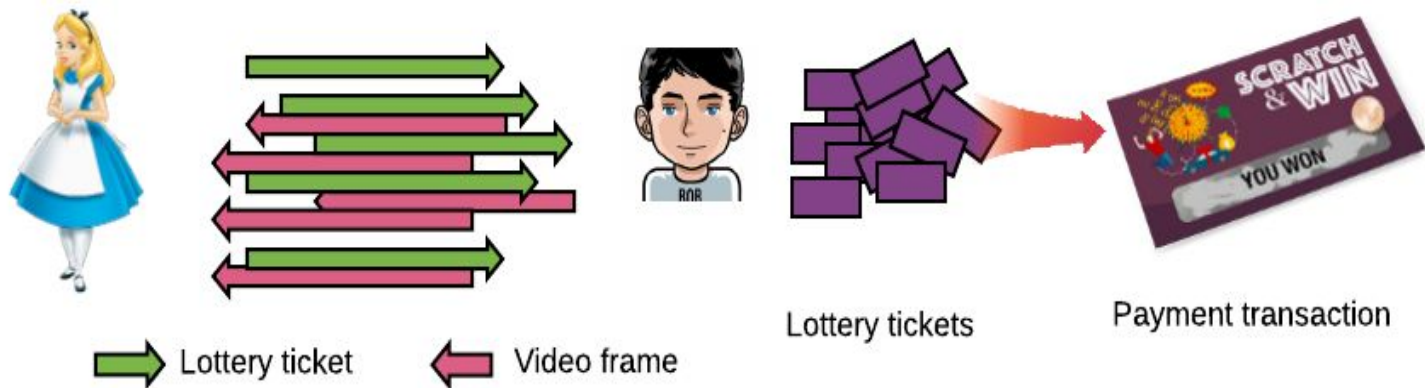


Issues

- Drive the system toward centralization.
 - Only wealthy parties can afford to be payment hubs.
- Hubs charge fees for relaying payments.
 - ***Fees are back!*** They may exceed the micropayment value itself.
- But, payment channels between long-term transacting parties (two parties) are still useful to handle micropayments.
- Currently payment networks are more geared towards enhancing scalability (i.e., transaction throughput rate) of cryptocurrencies.
 - For macropayments (or regular size payments).

Probabilistic Micropayments

- A solution to aggregate tiny payments.
- Dated back to Rivest [Rivest, 1997] and Wheeler [Wheeler, 1996].



Centralized Probabilistic Micropayments

- Early schemes were centralized.
- Involve a ***trusted bank*** to:
 - Authenticate users.
 - Hold users' accounts.
 - Authorize customers to issue lottery tickets.
 - Audit the lottery and manage payments.
- We will explore the scheme of [Rivest, 1997].
 - The original version that is based on an interactive coin tossing protocol.

Rivest's Scheme - Setup

- Beside creating accounts with the bank, the customer and merchant do the following:

- The customer creates a hash chain

$$x_0, x_1, x_2, \dots, x_n, \text{ where } x_i = H(x_{i+1}).$$

- The merchant creates a hash chain

$$y_0, y_1, y_2, \dots, y_n, \text{ where } y_i = H(y_{i+1}).$$

- The merchant sends the root y_0 (signed) to the customer.
- The customer sends the root x_0 concatenated with y_0 (signed) to the merchant.
 - This commits both parties to the hash chains they created.

Rivest's Scheme - Payments

- A customer pays a merchant at round i by sending them x_i .
- A micropayment wins if $x_i \bmod n = y_i \bmod n$
 - Where $n = 1/p$ (must be an integer).
- Upon winning, the merchant sends the committed chain roots, in addition to x_i and y_i , to the bank.
 - The bank verifies that the ticket is a winning valid one.
 - Checks the validity of the chain, the lottery outcome, signatures, etc.
 - Then it transfers currency from the customer's account to the merchant's account.

Drawbacks - Centralization!

- Increases the deployment cost.
 - Establishing relationships/accounts with bank.
- Limits the use of the payment service to systems with fully authenticated users.
- Trust and transparency issues!
 - If a distributed resource market, for example, uses it then the service is not decentralized anymore!
 - We are back to old P2P-based service models.

Decentralized Probabilistic Micropayments

- Utilize blockchain/cryptocurrencies to convert centralized schemes into distributed ones.
- Ingredients:
 - The bank is replaced with the miners.
 - Escrows/accounts are created on the blockchain.
 - Consensus/governance rules to manage escrows/accounts, claim/verify winning tickets, and punish cheaters.
- Will explore two systems:
 - **MICROPAY** [Pass et al., 2015],
 - **MicroCash** [Almashaqbeh et al., 2020].

MICROPAY1- Setup

- The customer creates an escrow with value X/p .
 - X is the expected value of a micropayment, and X/p is the value of a winning lottery ticket (i.e., total payment value).
 - This escrow can pay **only one** winning lottery ticket.
 - The escrow has its own public-private keypair.
 - The customer **knows** the private key of the escrow.
- So simply the customer creates a transaction transferring money to the escrow's address.

MICROPAY1 - Payment

- The merchant asks for a payment (or a lottery ticket) as follows:
 - Select a random number r_1 ,
 - Generate a commitment to r_1 called c (e.g., $c = \text{hash}(r_1)$).
 - Generate a public key pk_M .
 - Send (c, pk_M) signed to the customer.
- The customer replies as follows:
 - Select another random number r_2 ,
 - Send (r_2, c, pk_M) signed using the escrow's private key back to the merchant.
- So it is a two-round (interactive) lottery protocol.

MICROPAY1 - Lottery

- A ticket wins if:

$r1 \text{ XOR } r2$ has $\log(1/p)$ leading zero digits

(think about the XOR result in decimal).

- The merchant sends the lottery ticket (c, r1, r2, signature) to the miners.
 - This constitutes an unlocking script (in Bitcoin terms) to spend the escrow transaction.

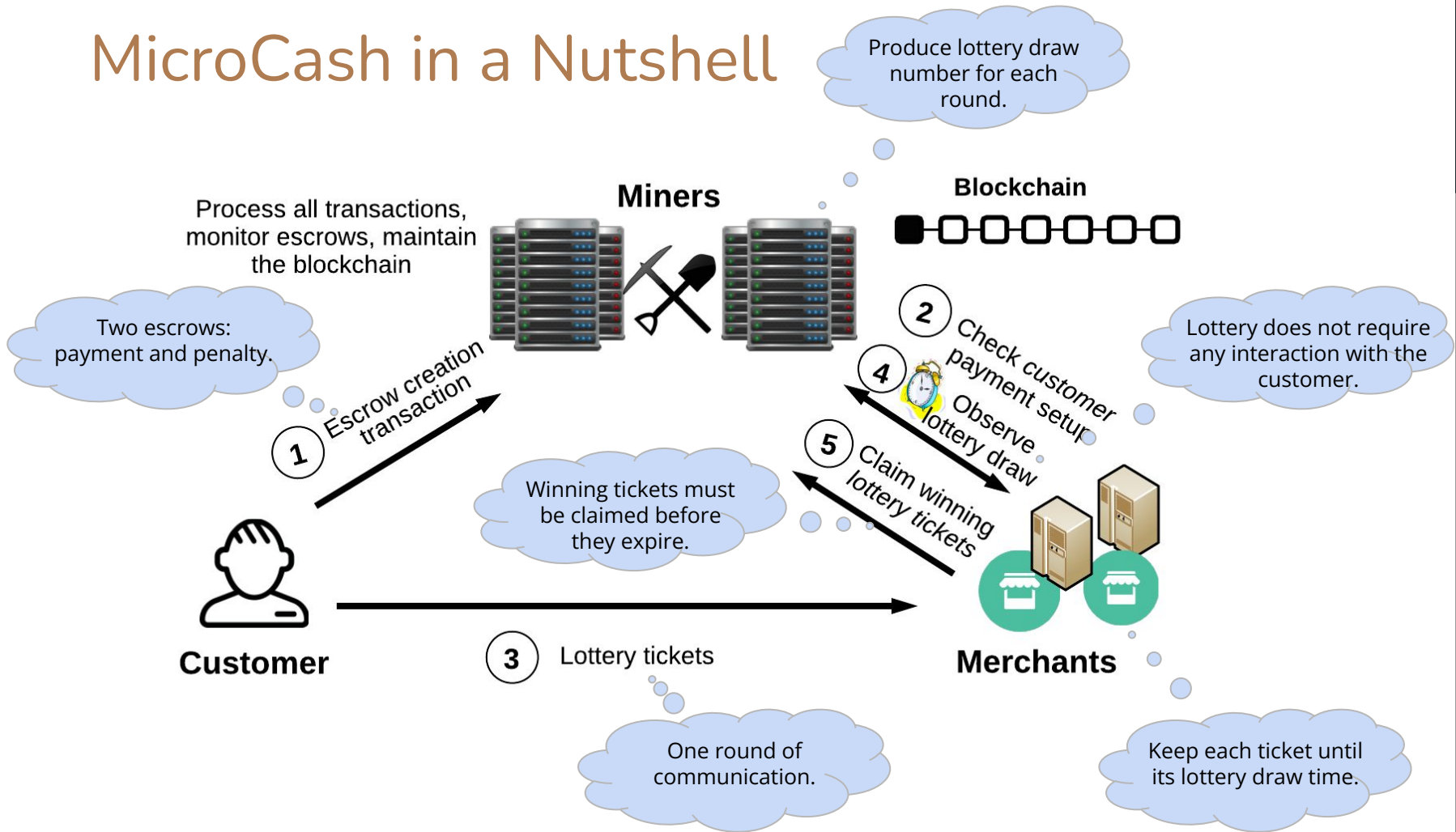
MICROPAY1 - Issues

- Several issues:
 - **Sequential** ticket issuance under the same escrow.
 - **Double spending**: issue the same ticket to several merchants.
 - *Why is this an issue?*
 - **Front running**: withdraw the escrow before a merchant claims its payment.
 - *Why is this an issue?*
 - Both are mitigated financially by having a penalty escrow.
 - However, the amount of this penalty is **not specified**.
 - **Interactive lottery**.
 - A non-interactive lottery was introduced but it is computationally heavy.
 - Chances of having **all tickets win or none** (psychological obstacle to use the system).

MicroCash

- The *first* decentralized probabilistic micropayment scheme that supports **concurrent micropayments**.
- The *first* to introduce a lottery with *exact win rate*.
 - Non-interactive lottery requiring only secure hashing (in the random oracle model).
- Compared to sequential micropayment schemes, it reduces the amount of data on the blockchain by around **50%**.
 - This is due to the fact an escrow can pay multiple winning tickets.
 - *So?*

MicroCash in a Nutshell

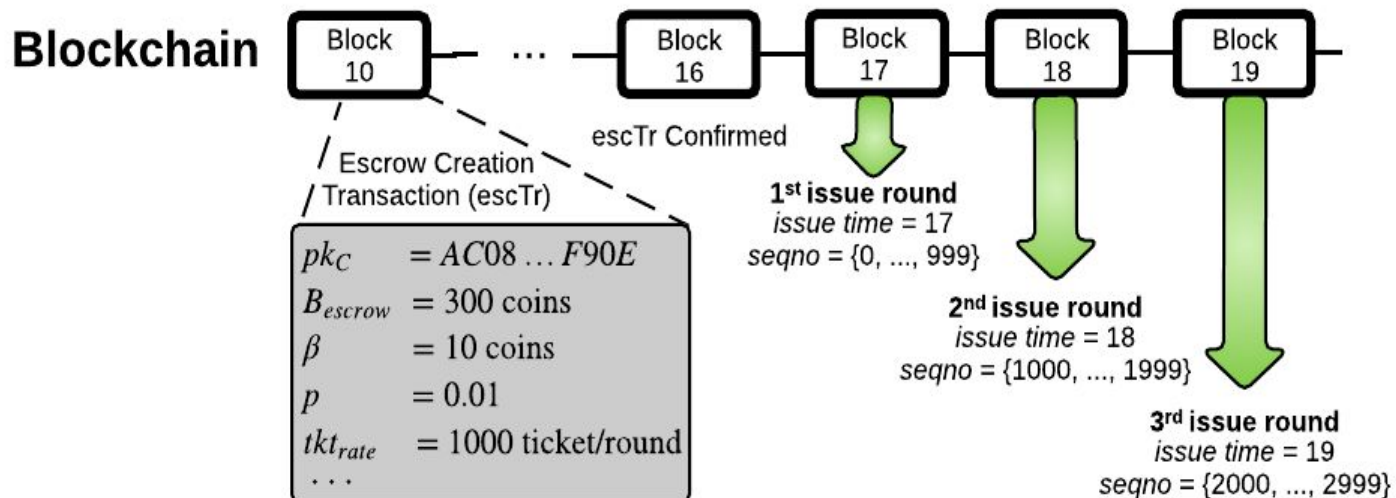


Lottery Ticket Issuance

- Each ticket is a simple structure consisting of:

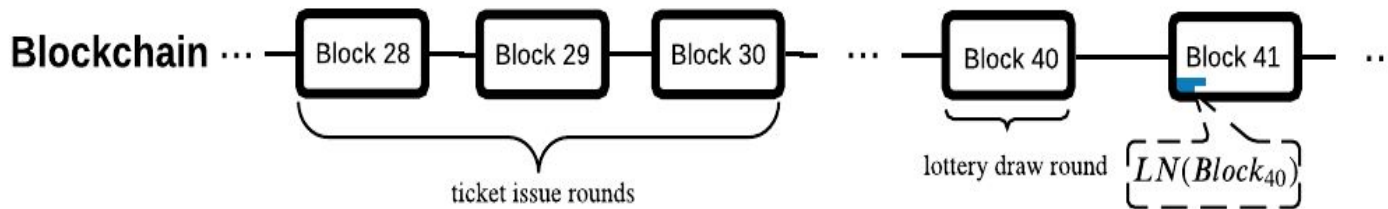
$$tkt_L = id_{esc} // index_M // seqno // \sigma_C$$

- Ticket issuance must follow a ticket issuing schedule.



The lottery Protocol

- Lightweight, non-interactive, and supports exact win rate.
 - Based on the blockchain view and requires only secure hashing.



tk_{tL}

id_{escrow}	= AC ... 0E
$seqno$	= 711
$index_M$	= 1
σ_C	= ...

$$h_1 = H(id_{escrow} || LN(Block_{40}))$$

$$\rightarrow \text{winning } seqno_1 = h_1 \% (tk_{t\text{rate}} \text{draw}_{len}) = 319$$

$$h_2 = H(h_1)$$

$$\rightarrow \text{winning } seqno_2 = h_2 \% (tk_{t\text{rate}} \text{draw}_{len}) = 711$$

...

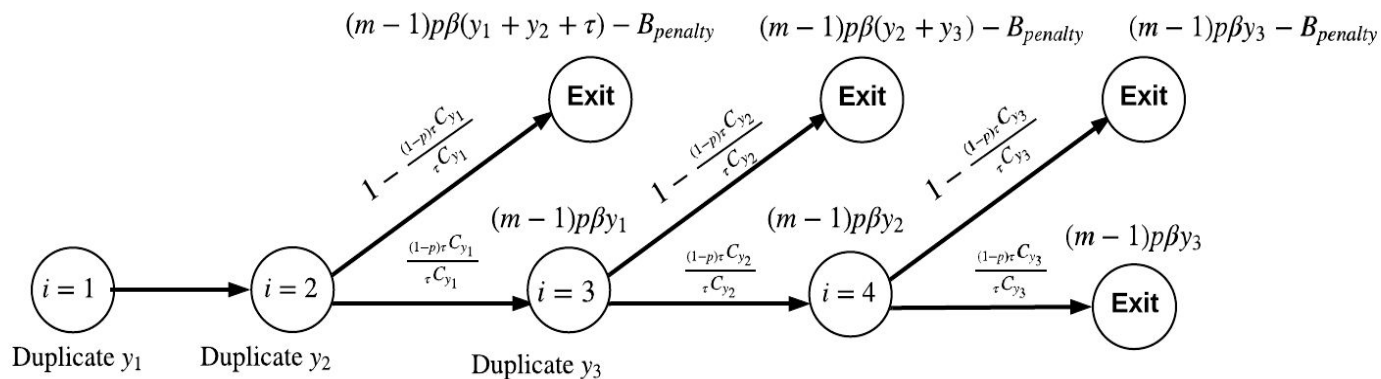
Winning Set
of Tickets



tk_{tL} is a
winning ticket

Penalty Escrow

- Used to defend against ticket duplication (or double spending).
 - Equals at least the additional utility a malicious customer obtains over an honest.



Theorem. For the game setup of MicroCash, issuing invalid or duplicated lottery tickets is less profitable in expectation than acting in an honest way if:

$$B_{penalty} > (m-1)p\beta\tau \left(\frac{1-p}{1 - \frac{1}{\tau C_{(1-p)\tau}}} + (1-p)(d-1) + r \right)$$

MicroCash - Issues

- Not fully compatible with any of the cryptocurrencies out there.
 - *Why?*
- To address double spending, the set of merchants that can be paid by using an escrow must be set in advance.
 - Needed for the lower bound analysis of the penalty deposit.
- Works in the random oracle model.

References

- [Rivest, 1997] Ronald Rivest. 1997. Electronic lottery tickets as micropayments. In International Conference on Financial Cryptography. Springer, 307–314.
- [Wheeler, 1996] David Wheeler. 1996. Transactions using bets. In International Workshop on Security Protocols. Springer, 89–92.
- [Pass et al., 2015] Pass, Rafael and Abhi Shelat. "Micropayments for decentralized currencies." In Proceedings of the 22nd ACM SIGSAC Conference on Computer and Communications Security, pp. 207-218. ACM, 2015.
- [Almashaqbeh et al., 2020] Ghada Almashaqbeh et al. "MicroCash: Practical Concurrent Processing of Micropayments." In Financial Cryptography, 2020.

